

POTENTIALIST SETS, INTENSIONS, AND NON-CLASSICALITY

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I thought I'd start with some **hot-off-the-press** philosophy:

*something is **infinite** if, taking it quantity by quantity, we can always take something **outside**. (Aristotle, .Phys. 207A7–15)*

This might be **old**, but similar intuitions have gained some currency in the philosophy of set theory **recently**.

*Set-theoretic potentialism is the view in the philosophy of mathematics that the universe of set theory is **never fully completed**, but rather **unfolds gradually** as parts of it increasingly **come into existence** or **become accessible** to us. [Hamkins and Linnebo, 2018, p. 1]*

- Alongside **potentialism** we have the idea of **intensions**.
- An **intension** is a collection that can **change members** (perhaps it's individuated by some **condition** or other).
- On the 'standard' view where we have an 'unchanging' universe, there isn't much room for intensions (since every intension just identifies a **single** extension).
- But under **potentialism** there's the possibility of having **interesting** intensions.
- How to **formalise these ideas**?

MAIN AIMS.

1. Present some of the **tools** that potentialists have used (and why they **won't work** in the intensional context).
2. Suggest a **different** option, and give some **preliminary** results about propositional logics.
3. Explain some **challenges** and **open questions** moving forward.

INTRODUCTION

POTENTIALISM AND MIRRORING

INTENSIONAL POTENTIALIST SYSTEMS

LIBERAL POTENTIALISM: BOOLEAN-VALUED

STRICT POTENTIALISM: HEYTING-VALUED

PROSPECTS FOR FURTHER RESULTS?

- Let's first discuss the **different kinds** of potentialism on offer.
- **Height potentialism** says that given some 'proper class' or other X , one **can** turn X into a set (you'll then have new "proper classes").
- **Rank potentialism** is the version of height potentialism where we can always form the **powerset** of a given domain.
- **Width potentialism** says that given some "powerset" $\mathcal{P}(x)$ or other, one **can** add new subsets not in $\mathcal{P}(x)$ (you'll then have a "new" powerset of x).
- **Forcing potentialism** is the version of width potentialism where we can always add a **generic** for any given partial order and family of dense sets.

- It's natural to formalise these ideas **modally**.
- I won't go into full details, but here's the **shape** of the formalisation.
- We add a new **modal operator** $\Diamond$ to the language.
- We define an appropriate **modal logic** (usually classical S4.2, maybe more) and give **axioms** for our potentialism in that logic.

- OK, but let's be clear: It would be **totally unreasonable** to expect mathematicians to adopt these modal theories.
- Thankfully there's a **translation** from the familiar language of set theory $\mathcal{L}_\in$ into our modal language.
- Again, suppressing details, one basically just replaces $\forall$ with $\Box\forall$ and $\exists$ with $\Diamond\exists$ (this is called the **potentialist translation**).

We can now prove some **nice theorems**:

THEOREM

[Linnebo, 2013] There is a modal theory of sets for **height potentialism** that interprets every axiom of ZFC under the potentialist translation.

THEOREM

[Scambler, MS] There is a modal theory of **height potentialism** and **width potentialism** that interprets every axiom of $\text{ZFC}^- +$ “Every set is countable” + Π_1^1 -Perfect Set Property.

“THEOREM”

[Barton, MS] There is a modal theory of **width potentialism** that interprets $\text{ZFC}^- +$ “Every set is countable” + Π_1^1 -Perfect Set Property.

- This is great for claims about **sets**.
- We can read the “normal” mathematical quantifiers as **implicit** modal quantifiers.
- But there’s a couple of **problems**.
- **Problem 1.** We need **classical logic**, and this is a little controversial (stay tuned).
- **Problem 2.** For intensions, the proofs depend on the following **stability axioms**:
 - (**Positive stability.**) $x \in y \rightarrow \Box x \in y$
 - (**Negative stability.**) $x \notin y \rightarrow \Box x \notin y$
- **Neither** of these looks very plausible for intensions.
- Consider:
 - The intension given by the condition “ x is constructible” under **height potentialism**.
 - The intension given by the condition “ x is countable” under **width potentialism**.

- **Two** ways of finding stuff out about a subject matter:
 1. Write down some good **axioms**.
 2. Study a bunch of **structures**.
- Option 1 looks **dead** using current technology (at least for intensions).
- Option 2 shows some **promise**.

DEFINITION.

[Hamkins and Linnebo, 2018] A **potentialist system** (for sets) is a pair $\mathbb{S} = (S, \leq_{\mathbb{S}})$, where S is a collection of structures of a certain kind and $\leq_{\mathbb{S}}$ is a **refinement** of the substructure relation.

- We'll modify this approach to fit **intensions**.
- First let's see some **examples**.
- Fix some **countable transitive model** $M \models \text{ZFC}$.

DEFINITION.

The Rank Potentialist System is the following potentialist system:

- (I) $S = \{V_\beta^M \mid \beta \in On^M\}$
- (II) $V_\alpha^M \leq_S V_\beta^M$ iff $\alpha \leq \beta$.

DEFINITION.

The Woodin Generic Multiverse Potentialist System is the following potentialist system:

- (I) $S = \{M[G] \mid \text{"}G \text{ is } M\text{-generic"}\}$
- (II) $M[G] \leq_S M[H]$ iff $M[H]$ is a forcing extension of $M[G]$.

DEFINITION

Let $M[G]$ be a $Col(\omega, < Ord)$ class forcing extension of M . Then *the Steel Generic Multiverse Potentialist System* is the following potentialist system:

1. $S = \{x \mid "x = M[H] \text{ and there is an } \alpha \in M \text{ such that } H \in M[G \restriction \alpha]" \}$
2. $M[H_0] \leq_S M[H_1]$ iff $M[H_1]$ is a forcing extension of $M[H_0]$.

DEFINITION.

The Countable Transitive Potentialist System is the following potentialist system:

- (I) $S = \{N \mid \text{"}N \text{ is a countable transitive model of ZFC containing } M\text{"}\}$
- (II) $N \leq_S N'$ iff $N \subseteq N'$.

DEFINITION.

The Actualist (Potentialist) System is the following potentialist system:

- (I) $S = \{M\}$
- (II) $M \leq_S M$ iff $M = M$

- We now want to add in **intensions**.
- Start by considering the **full-domain** of $\mathbb{S}$,
 $fulldom(\mathbb{S}) = \bigcup_{W \in \mathbb{S}} W$.
- We'll want to allow **new** intensions to emerge, so we'll handle this by using partial functions from $\mathbb{S}$ to $fulldom(\mathbb{S})$.
- **Choice point:** I **don't** want intensions to pop out of existence, so we'll add the following requirement.

DEFINITION.

A partial function $f : \mathbb{S} \rightarrow \mathcal{P}(fulldom(\mathbb{S}))$ is *upwards defined on $\mathbb{S}$* if whenever $f(W)$ is defined, $f(W')$ is also defined for all $W' \geq_{\mathbb{S}} W$.

- **Choice point:** We also don't want functions to **reach beyond** the domains over which they are define. So we'll add the following:

DEFINITION.

An upwards defined partial function $f : S \rightarrow \mathcal{P}(\bigcup_{W \in S} W)$ is *S-bounded* iff $f(W) \subseteq W$ for every $W \in \text{dom}(f)$.

DEFINITION.

An *intensional potentialist system* is a triple $\mathbb{S} = (S, \leq_{\mathbb{S}}, \mathcal{X})$, where $(S, \leq_{\mathbb{S}})$ is a **potentialist system** and $\mathcal{X}$ is a collection of upwards-closed $\mathbb{S}$ -bounded partial functions for $(S, \leq_{\mathbb{S}})$.

An **important** example:

DEFINITION

The *definabilist intensional system over $\mathbb{S}$* is the system that has as intensions all definable intensions at every given world.

(More formally, given any finite list of parameters $\vec{x}$ and formula $\phi(\vec{x}, y)$, where there is a world $W \in \mathbb{S}$ containing each parameter $\bar{x}$ in $\vec{x}$, there is a $\mathbb{S}$ -bounded partial function $f : \mathbb{S} \rightarrow \mathcal{P}(\text{fulldom}(\mathbb{S}))$ such that $f : W \mapsto \{y | W \models \phi(\vec{x}, y)\}.$)

- Some **important** examples.
- **Def.** We say that f is *constant* if $f(W) = f(W')$ for any worlds W and W' on which f is defined.
- **Def.** We say that f is *set-like* if there is an $x \in \text{fulldom}(\mathbb{S})$ such that wherever $f(W)$ is defined, $f(W) = x$.
- **Def.** *Positively stable* (or ‘*monotone*’ or ‘*monotonic*’) iff whenever we have $x \in f(W)$, then for every $W' \geq_{\mathbb{S}} W$, we have $x \in f(W')$.
- **Def.** *Negatively stable* iff whenever we have $x \in f(W)$, then for every $W' \leq_{\mathbb{S}} W$, we have $x \in f(W')$.
- **Def.** *Stable* iff f is both positively stable and negatively stable.
- **Def.** *Capricious* iff for some set $x \in \text{fulldom}(\mathbb{S})$ and every world W with $x \in W$, there are $W_1, W_2 \geq_{\mathbb{S}} W$ such that $x \notin f(W_1)$ and $x \in f(W_2)$.

- **Def.** We say that f is the *stabilisation* of g iff whenever there is an $x \in \text{fulldom}(\mathbb{S})$ and W such that $x \in g(W)$ then $x \in f(W')$ for every world at which x exists.
- **Def.** We define f' , the *component-wise complement of f* as the partial function $f'(W) = W/f(W)$ (and is undefined wherever f is).
- **More** classes besides (including **weakenings**) of these notions.

- What we now want is some way of assigning **truth values** to set-theoretic claims.
- To begin with, let's just treat the **propositional case**.
- Here we treat every atomic and quantified sentence as a **propositional variable** and can then talk about how we assign truth-values to these claims using **compositional clauses**.

CORE INTUITION.

- Each $\mathbb{S}$ gives us a **partial order** of the worlds under $\leq_{\mathbb{S}}$.
- We can use this partial order to consider particular **algebras**.
- We then use **algebraic semantics** to get truth-values for set-theoretic sentences.

Important here is the following distinction:

LIBERAL POTENTIALISM

Liberal potentialists regard the modal truths as **unproblematic**. In particular, there are modal truths about generative processes **in their entirety**, including those that cannot be completed...

STRICT POTENTIALISM

strict potentialism...goes beyond the liberal view by requiring, not only that every object be generated at some stage of a process, but also that every truth be “**made true**” at some stage. [Linnebo and Shapiro, 2019]

- How should we assign a **semantics**?
- **Slogan:** Ask not **if** a sentence is true, but **when**.
- For **liberal potentialism** consider the following:

DEFINITION

- Given an intensional potentialist system $\mathbb{S}$, we let the **Boolean algebra for $\mathbb{S}$** (or $\mathbb{B}(\mathbb{S})$) be simply the powerset of the collection of all worlds of $\mathbb{S}$.
- Algebraic operations handled exactly as in the powerset.

- We now assign elements of $\mathbb{B}(\mathcal{S})$ via:
- **Propositional variables** (e.g. atomics and quantified statements): $\llbracket \phi \rrbracket = \{W \mid W \models \phi\}$.
- **Compositional clauses** handled by the algebraic operations.

FACT.

Because $\mathbb{B}(\mathbb{S})$ is **always** a Boolean-algebra, the propositional logic for **any** potentialist system is **classical logic**.

- What about the **strict** potentialist?
- The strict potentialist thinks that claims about sets and classes are **made true** as we move through $\mathbb{S}$.
- This has more in keeping with a **intuitionistic/constructive** approach.

- We now modify our algebra we want to consider to $Up(\mathbb{S})$:
The collection of all upsets of $\mathbb{S}$.
- **Fact:** This is a **Heyting** algebra (with pseudocomplements and relative pseudocomplements defined as normal).
- Again, we assign semantics via the **algebraic operations**.
- But here we assign the value of a propositional variable $\llbracket \phi \rrbracket$ according to whether ϕ is satisfied at W **and** ϕ is satisfied at **every** $W' \geq_{\mathbb{S}} W$.
- **Slogan:** ϕ is true **now** and **forevermore**!

- The calibration of the propositional logic then **depends** on the structure of $\mathbb{S}$.
- In particular, when it is **linear** we get Gödel-Dummett logic: all classical tautologies of the form $(\psi_0 \rightarrow \psi_1) \vee (\psi_1 \rightarrow \psi_0)$ added.
- This means that the propositional logic we get is **stronger** for just the **rank potentialist** system as opposed to **any** width potentialism.
- But in all cases we see failures of LEM.
- e.g. “ x is constructible or x is not constructible” under rank potentialism.
- e.g. Allowing me a parameter $\bar{x}$ for a set x , “ $\bar{x}$ is countable or $\bar{x}$ is uncountable” under any kind of forcing potentialism.

- Some **conclusions**:
- The propositional logic you get depends on whether you're a **strict** or **liberal** potentialist (suitably formalised).
- **Unsurprising**. **Liberal potentialists** get **classical logic**, **strict potentialists** get something **more constructive**.
- **More surprising**. The **kind of potentialism** can affect the **amount of classicality** you get.
- In particular, **rank potentialism** is **substantially more classical** than any kind of **forcing potentialism**.

- **Key open question.** How to generalise this to **first-order** theories (possibly two-sorted to allow for intensions)?
- **Objective.** Provide a semantics for first-order (not just propositional) claims, and examine what goes on.
- This looks **hard** as **many** classes are **not monotonic**.
- **Ideas:**
 - **Algebra-valued** models?
 - **Iemhoff-Paßmann** models?
 - **Lubarsky** models?
 - **Presheaves**?

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